



2. Триангульовані категорії. Навколо похідних категорій

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Властивості $\mathbf{K}(\mathcal{C})$ як триангульованої категорії

Proposition 1.4.4. *The collection of distinguished triangles in $\mathbf{K}(\mathcal{C})$ satisfies the following properties, (TR 0)–(TR 5).*

(TR 0) *A triangle isomorphic to a distinguished triangle is distinguished.*

(TR 1) *For any $X \in \text{Ob}(\mathbf{K}(\mathcal{C}))$, $X \xrightarrow{\text{id}_X} X \longrightarrow 0 \longrightarrow X[1]$ is a distinguished triangle.*

(TR 2) *Any $f : X \rightarrow Y$ in $\mathbf{K}(\mathcal{C})$ can be embedded in a distinguished triangle $X \xrightarrow{f} Y \rightarrow Z \rightarrow X[1]$.*

(TR 3) *$X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} X[1]$ is a distinguished triangle if and only if $Y \xrightarrow{g} Z \xrightarrow{h} X[1] \xrightarrow{-f[1]} Y[1]$ is a distinguished triangle.*

(TR 4) *Given two distinguished triangles $X \xrightarrow{f} Y \rightarrow Z \rightarrow X[1]$ and $X' \xrightarrow{f'} Y' \rightarrow Z' \rightarrow X'[1]$, a commutative diagram*

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow u & & \downarrow v \\ X' & \xrightarrow{f'} & Y' \end{array}$$

can be embedded in a morphism of triangles (not necessarily unique).

(TR 5) (octahedral axiom). *Suppose given distinguished triangles:*

$$\begin{array}{ccccccc} X & \xrightarrow{f} & Y & \longrightarrow & Z' & \longrightarrow & X[1] , \\ Y & \xrightarrow{g} & Z & \longrightarrow & X' & \longrightarrow & Y[1] , \\ X & \xrightarrow{g \circ f} & Z & \longrightarrow & Y' & \longrightarrow & X[1] , \end{array}$$

then there exists a distinguished triangle

$$Z' \rightarrow Y' \rightarrow X' \rightarrow Z'[1]$$

such that the following diagram is commutative:

$$\begin{array}{ccccccc} X & \xrightarrow{f} & Y & \longrightarrow & Z' & \longrightarrow & X[1] \\ \text{id}_X \downarrow & & \downarrow g & & \downarrow & & \downarrow \text{id}_{X[1]} \\ X & \xrightarrow{g \circ f} & Z & \longrightarrow & Y' & \longrightarrow & X[1] \\ f \downarrow & & \downarrow \text{id}_Z & & \downarrow & & \downarrow f[1] \\ Y & \xrightarrow{g} & Z & \longrightarrow & X' & \longrightarrow & Y[1] \\ \downarrow & & \downarrow & & \downarrow \text{id}_{X'} & & \downarrow \\ Z' & \longrightarrow & Y' & \longrightarrow & X' & \longrightarrow & Z'[1] . \end{array}$$

Доведення властивостей $\mathbf{K}(\mathcal{C})$ як трианг. категорії

Proof. The properties (TR 0) and (TR 2) are obvious, and (TR 3) follows from Lemma 1.4.2.

Since the mapping cone of $f : 0 \rightarrow X$ is X , the triangle $0 \rightarrow X \xrightarrow{\text{id}_X} X \rightarrow 0[1]$ is distinguished. Applying (TR 3) we get (TR 1). Let us prove (TR 4). We may assume that $X \xrightarrow{f} Y \rightarrow Z \rightarrow X[1]$ and $X' \xrightarrow{f'} Y' \rightarrow Z' \rightarrow X'[1]$ are $X \xrightarrow{f} Y \xrightarrow{\alpha(f)} M(f) \xrightarrow{\beta(f)} X[1]$ and $X' \xrightarrow{f'} Y' \xrightarrow{\alpha(f')} M(f') \xrightarrow{\beta(f')} X'[1]$, respectively. We shall construct a morphism $w : M(f) \rightarrow M(f')$ such that:

$$(1.4.4) \quad \begin{cases} w \circ \alpha(f) = \alpha(f') \circ v , \\ u[1] \circ \beta(f) = \beta(f') \circ w . \end{cases}$$

By the definition of $\mathbf{K}(\mathcal{C})$, there exists $s^n : X^n \rightarrow Y'^{n-1}$ such that $v^n \circ f^n - f'^n \circ u^n = s^{n+1} \circ d_X^n + d_{Y'}^{n-1} \circ s^n$. We define $w^n : M(f)^n = X^{n+1} \oplus Y^n \rightarrow M(f')^n = X'^{n+1} \oplus Y'^n$ by:

$$w^n = \begin{pmatrix} u^{n+1} & 0 \\ s^{n+1} & v^n \end{pmatrix} .$$

Then a direct calculation shows that w is a morphism of complexes and satisfies (1.4.4).

$$w = \begin{pmatrix} u[1] & 0 \\ s \circ \sigma^{-1} & v \end{pmatrix} .$$

Let us prove (TR 5). We may assume $Z' = M(f)$, $X' = M(g)$ and $Y' = M(g \circ f)$. Let us define $u: Z' \rightarrow Y'$ and $v: Y' \rightarrow X'$ by:

$$u^n: X^{n+1} \oplus Y^n \rightarrow X^{n+1} \oplus Z^n, \quad u = \begin{pmatrix} \text{id}_{X^{n+1}} & 0 \\ 0 & g^n \end{pmatrix},$$

$$v^n: X^{n+1} \oplus Z^n \rightarrow Y^{n+1} \oplus Z^n, \quad v = \begin{pmatrix} f^{n+1} & 0 \\ 0 & \text{id}_{Z^n} \end{pmatrix}.$$

$$\begin{array}{ccccccc}
 X & \xrightarrow{f} & Y & \xrightarrow{\alpha(f)} & Z' & \xrightarrow{\beta(f)} & X[1] \\
 \text{id}_X \downarrow & & \downarrow g & & \downarrow u & & \downarrow \text{id}_{X[1]} \\
 X & \xrightarrow{g \circ f} & Z & \xrightarrow{\alpha(g \circ f)} & Y' & \xrightarrow{\beta(g \circ f)} & X[1] \\
 f \downarrow & & \downarrow \text{id}_Z & & \downarrow v & & \downarrow f[1] \\
 Y & \xrightarrow{g} & Z & \xrightarrow{\alpha(g)} & X' & \xrightarrow{\beta(g)} & Y[1] \\
 \alpha(f) \downarrow & & \downarrow \alpha(g \circ f) & & \downarrow \text{id}_{X'} & & \downarrow \alpha(f)[1] \\
 Z' & \xrightarrow{u} & Y' & \xrightarrow{v} & X' & \longrightarrow & Z'[1] .
 \end{array}$$

We define $w: X' \rightarrow Z'[1]$ as the composite $X' \rightarrow Y[1] \rightarrow Z'[1]$. Then the diagram in (TR 5) is commutative, and it is enough to show that $Z' \xrightarrow{u} Y' \xrightarrow{v} X' \xrightarrow{w} Z'[1]$ is a distinguished triangle. For that purpose we shall construct an isomorphism $\phi: M(u) \rightarrow X'$ and its inverse $\psi: X' \rightarrow M(u)$ such that $\phi \circ \alpha(u) = v$ and $\beta(u) \circ \psi = w$. We have:

$$M(u)^n = M(f)^{n+1} \oplus M(g \circ f)^n = X^{n+2} \oplus Y^{n+1} \oplus X^{n+1} \oplus Z^n$$

and $X'^n = M(g)^n = Y^{n+1} \oplus Z^n$. We define ϕ and ψ by:

$$\phi^n = \begin{pmatrix} 0 & \text{id}_{Y^{n+1}} & f^{n+1} & 0 \\ 0 & 0 & 0 & \text{id}_{Z^n} \end{pmatrix}, \quad \psi^n = \begin{pmatrix} 0 & 0 \\ \text{id}_{Y^{n+1}} & 0 \\ 0 & 0 \\ 0 & \text{id}_{X^{n+1}} \end{pmatrix}.$$

Then one checks easily that ϕ and ψ are morphisms of complexes and $\phi \circ \alpha(u) = v$, $\beta(u) \circ \psi = w$.

$$\begin{array}{ccccccc} X & \xrightarrow{f} & Y & \longrightarrow & Z' & \longrightarrow & X[1] \\ \text{id}_X \downarrow & & \downarrow g & & \downarrow u & & \downarrow \text{id}_{X[1]} \\ X & \xrightarrow{g \circ f} & Z & \longrightarrow & Y' & \longrightarrow & X[1] \\ f \downarrow & & \downarrow \text{id}_Z & & \downarrow v & & \downarrow f[1] \\ Y & \xrightarrow{g} & Z & \longrightarrow & X' & \longrightarrow & Y[1] \\ \downarrow & & \downarrow & & \downarrow \text{id}_{X'} & \searrow w & \downarrow \\ Z' & \xrightarrow{u} & Y' & \xrightarrow{v} & X' & \xrightarrow{w} & Z'[1] \end{array}.$$

Октаэдр

We have $\phi \circ \psi = \text{id}_{X'}$. If we define:

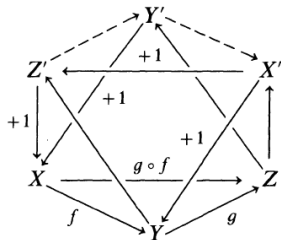
$$s^n : M(u)^n \rightarrow M(u)^{n-1} , \quad s^n = \begin{pmatrix} 0 & 0 & \text{id}_{X^{n+1}} & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

then:

$$(\text{id}_{M(u)} - \psi \circ \phi)^n = s^{n+1} \circ d_{M(u)}^n + d_{M(u)}^{n-1} \circ s^n .$$

Hence $\psi \circ \phi$ equals $\text{id}_{M(u)}$ in $\mathbf{K}(\mathcal{C})$. \square

Remark 1.4.5. Property (TR 5) may be visualized by the following octahedral diagram:



Триангульована категорія

Let \mathcal{C} be an additive category, together with an automorphism $T : \mathcal{C} \rightarrow \mathcal{C}$. We write sometimes $[1]$ for T and $[k]$ for T^k , (i.e. $X[1]$ for $T(X)$, or $f[1]$ for $T(f)$).

A triangle in \mathcal{C} is a sequence of morphisms

$$X \rightarrow Y \rightarrow Z \rightarrow T(X) .$$

Definition 1.5.1. *A triangulated category \mathcal{C} consists of the following data and rules.*

(1.5.1) *An additive category \mathcal{C} together with an automorphism $T : \mathcal{C} \rightarrow \mathcal{C}$,*

(1.5.2) *a family of triangles, called distinguished triangles.*

These data satisfy the axioms (TR 0)–(TR 5) of Proposition 1.4.4 when setting $X[1] = T(X)$.

Let (\mathcal{C}, T) and (\mathcal{C}', T') be two triangulated categories. We say that an additive functor F from \mathcal{C} to \mathcal{C}' is a functor of triangulated categories if $F \circ T \simeq T' \circ F$, and F sends distinguished triangles of \mathcal{C} into distinguished triangles of \mathcal{C}' .

Clearly, for an additive category \mathcal{C} , $\mathbf{K}(\mathcal{C})$ is a triangulated category.

Когомологічний функтор

\mathcal{C} – триангульована категорія, $\mathcal{A} = \mathbb{k}\text{-mod}$.

Definition 1.5.2. An additive functor $F : \mathcal{C} \rightarrow \mathcal{A}$ is called a cohomological functor if for any distinguished triangle $X \rightarrow Y \rightarrow Z \rightarrow T(X)$, the sequence $F(X) \rightarrow F(Y) \rightarrow F(Z)$ is exact.

For a cohomological functor F , we write F^k for $F \circ T^k$. Then for any distinguished triangle $X \rightarrow Y \rightarrow Z \rightarrow T(X)$ we obtain a long exact sequence:

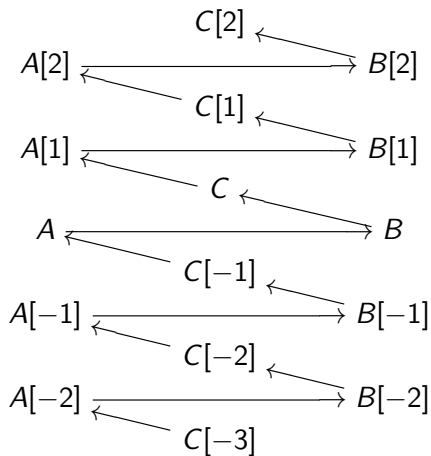
$$(1.5.3) \quad \cdots \rightarrow F^{k-1}(Z) \rightarrow F^k(X) \rightarrow F^k(Y) \rightarrow F^k(Z) \rightarrow F^{k+1}(X) \rightarrow \cdots .$$

Трикутник $A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} A[1]$ - скорочений запис послідовності

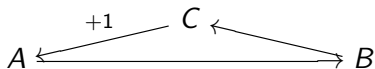
$$\cdots \xrightarrow{g[-1]} C[-1] \xrightarrow{h[-1]} A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{h} A[1] \xrightarrow{f[1]} B[1] \xrightarrow{g[1]} C[1] \xrightarrow{h[1]} \cdots$$

яку можна візуалізувати як спіраль

Трикутник



що проектується на трикутник



Proposition 1.5.3. (i) If $X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow T(X)$ is a distinguished triangle, then $g \circ f = 0$.

(ii) For any $W \in \text{Ob}(\mathcal{C})$, $\text{Hom}_{\mathcal{C}}(W, \cdot)$ and $\text{Hom}_{\mathcal{C}}(\cdot, W)$ are cohomological functors.

Proof. (i) By (TR 1), $X \xrightarrow{\text{id}_X} X \rightarrow 0 \rightarrow T(X)$ is a distinguished triangle. Therefore by (TR 4) there is a morphism $\phi: 0 \rightarrow Z$ which makes the following diagram commutative:

$$\begin{array}{ccccccc}
 X & \longrightarrow & X & \longrightarrow & 0 & \longrightarrow & T(X) \\
 \text{id}_X \downarrow & & f \downarrow & & \vdots \downarrow \phi & & \downarrow \\
 X & \xrightarrow{f} & Y & \xrightarrow{g} & Z & \longrightarrow & T(X) .
 \end{array}$$

Hence $g \circ f = \phi \circ 0 = 0$.

(ii) Let $X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow T(X)$ be a distinguished triangle. In order to show that $\text{Hom}_{\mathcal{C}}(W, \cdot)$ is a cohomological functor, it is enough to show that, for any $\phi \in \text{Hom}_{\mathcal{C}}(W, Y)$ with $g \circ \phi = 0$, we can find $\psi \in \text{Hom}_{\mathcal{C}}(W, X)$, with $\phi = f \circ \psi$. This follows from (TR 1), (TR 3) and (TR 4) which imply that the dotted arrow below can be completed:

$$\begin{array}{ccccccc}
 W & \xrightarrow{\text{id}_W} & W & \longrightarrow & 0 & \longrightarrow & T(W) \\
 \vdots \downarrow \psi & & \downarrow \phi & & \downarrow & & \\
 X & \xrightarrow{f} & Y & \xrightarrow{g} & Z & \longrightarrow & T(X) .
 \end{array}$$

The proof that $\text{Hom}_{\mathcal{C}}(\cdot, W)$ is a cohomological functor is similar. \square

5 ізоморфізмів

Corollary 1.5.5. *Let*

$$\begin{array}{ccccccc} X & \longrightarrow & Y & \longrightarrow & Z & \longrightarrow & T(X) \\ \phi \downarrow & & \psi \downarrow & & \theta \downarrow & & T(\phi) \downarrow \\ X' & \longrightarrow & Y' & \longrightarrow & Z' & \longrightarrow & T(X') \end{array}$$

be a morphism of distinguished triangles. If ϕ and ψ are isomorphisms, then so is θ .

Proof. For any $W \in \text{Ob}(\mathcal{C})$, let us apply the functor $\text{Hom}_{\mathcal{C}}(W, \cdot)$ to the above diagram. We obtain a commutative diagram whose rows are exact. Since $\text{Hom}_{\mathcal{C}}(W, \phi)$ and $\text{Hom}_{\mathcal{C}}(W, \psi)$ are isomorphisms, as well as $\text{Hom}_{\mathcal{C}}(W, T(\phi))$ and $\text{Hom}_{\mathcal{C}}(W, T(\psi))$, we obtain that $\text{Hom}_{\mathcal{C}}(W, \theta)$ is an isomorphism by Exercise I.8.

(5-лема). За лемою Йонеда θ - ізоморфізм.

$\mathcal{C} = \mathbb{k}\text{-mod}$.

Proposition 1.5.6. *Let \mathcal{C} be an abelian category. Then the functor $H^0(\cdot) : \mathbf{K}(\mathcal{C}) \rightarrow \mathcal{C}$ is a cohomological functor.*

Proof. It is enough to show that if $f : X \rightarrow Y$ is a morphism in $\mathbf{C}(\mathcal{C})$, then the sequence

$$H^0(Y) \rightarrow H^0(M(f)) \rightarrow H^0(X[1])$$

is exact.

Since $0 \rightarrow Y \rightarrow M(f) \rightarrow X[1] \rightarrow 0$ is an exact sequence in $\mathbf{C}(\mathcal{C})$, the result follows from Proposition 1.3.6. \square

Definition 1.5.7. *Let \mathcal{C} be an abelian category and let $f : X \rightarrow Y$ be a morphism in $\mathbf{K}(\mathcal{C})$. One says that f is a quasi-isomorphism (qis for short) if $H^n(f)$ is an isomorphism for each n .*

Hence f is a qis if and only if $H^n(M(f)) = 0$ for each n . If f is a qis, one writes $X \xrightarrow[\text{qis}]{} Y$, for short.

Notations 1.5.8. Let \mathcal{C} be a triangulated category. In the subsequent sections we shall often write $X \longrightarrow Y \longrightarrow Z \xrightarrow{+1}$ instead of $X \rightarrow Y \rightarrow Z \rightarrow T(X)$, to denote a distinguished triangle.

Let \mathcal{C} be a category, and let S be a family of morphisms in \mathcal{C} .

Definition 1.6.1. One says that S is a multiplicative system if it satisfies (S1)–(S4) below.

(S1) For any $X \in \text{Ob}(\mathcal{C})$, $\text{id}_X \in S$.

(S2) For any pair (f, g) of S such that the composition $g \circ f$ exists, $g \circ f \in S$.

(S3) Any diagram:

$$\begin{array}{ccc} & & Z \\ & & \downarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

with $g \in S$, may be completed to a commutative diagram:

$$\begin{array}{ccc} W & \longrightarrow & Z \\ \downarrow h & & \downarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

with $h \in S$. Ditto with all the arrows reversed.

(S4) If f and g belong to $\text{Hom}_{\mathcal{C}}(X, Y)$, the following conditions are equivalent:

- (i) there exists $t : Y \rightarrow Y'$, $t \in S$, such that $t \circ f = t \circ g$,
- (ii) there exists $s : X' \rightarrow X$, $s \in S$, such that $f \circ s = g \circ s$.

Локалізація

Definition 1.6.2. Let \mathcal{C} be a category, S a multiplicative system. The category \mathcal{C}_S , called the localization of \mathcal{C} by S , is defined by:

$$(1.6.1) \quad \text{Ob}(\mathcal{C}_S) = \text{Ob}(\mathcal{C}) ,$$

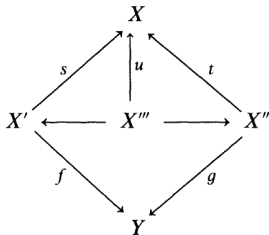
(1.6.2) for any pair (X, Y) of $\text{Ob}(\mathcal{C})$,

$$\text{Hom}_{\mathcal{C}_S}(X, Y) = \{(X', s, f); X' \in \text{Ob}(\mathcal{C}), s: X' \rightarrow X, f: X' \rightarrow Y, s \in S\} / \mathcal{R}$$

where \mathcal{R} is the following equivalence relation:

$$(X', s, f) \mathcal{R} (X'', t, g)$$

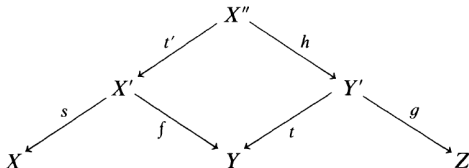
iff there exists a commutative diagram



with $u \in S$.

Композиція в локалізації

The composition of $(X', s, f) \in \text{Hom}_{\mathcal{C}_S}(X, Y)$ and $(Y', t, g) \in \text{Hom}_{\mathcal{C}_S}(Y, Z)$ is defined as follows. We use (S3) to find a commutative diagram:



with $t' \in S$, and we set:

$$(Y', t, g) \circ (X', s, f) = (X'', s \circ t', g \circ h) .$$

One sees easily, using the axioms (S1)–(S4), that \mathcal{C}_S is a category.

We shall denote by Q the functor:

$$Q: \mathcal{C} \rightarrow \mathcal{C}_S$$

defined by $Q(X) = X$ for $X \in \text{Ob}(\mathcal{C})$, and $Q(f) = (X, \text{id}_X, f)$ for $f \in \text{Hom}_{\mathcal{C}}(X, Y)$.

Proposition 1.6.3. (i) For $s \in S$, $Q(s)$ is an isomorphism in \mathcal{C}_S .

(ii) Let \mathcal{C}' be another category, $F: \mathcal{C} \rightarrow \mathcal{C}'$ be a functor such that $F(s)$ is an isomorphism for all $s \in S$. Then F factors uniquely through Q .

Нульова система

Definition 1.6.6. Let \mathcal{C} be a triangulated category, and let \mathcal{N} be a subfamily of $\text{Ob}(\mathcal{C})$. One says that \mathcal{N} is a null system if it satisfies (N 1)–(N3) below.

(N 1) $0 \in \mathcal{N}$,

(N 2) $X \in \mathcal{N}$ if and only if $X[1] \in \mathcal{N}$,

(N 3) If $X \rightarrow Y \rightarrow Z \rightarrow X[1]$ is a distinguished triangle, and $X \in \mathcal{N}$, $Y \in \mathcal{N}$, then $Z \in \mathcal{N}$.

Now we set:

$$(1.6.4) \quad \left\{ \begin{array}{l} S(\mathcal{N}) = \{f : X \rightarrow Y; f \text{ is embedded into a distinguished} \\ \text{triangle } X \xrightarrow{f} Y \rightarrow Z \rightarrow X[1], \text{ with } Z \in \mathcal{N}\} . \end{array} \right.$$

Proposition 1.6.7. *Assume \mathcal{N} is a null system. Then $S(\mathcal{N})$ is a multiplicative system.*

Proof. The property (S 1) is deduced from (N 1) and (TR 1). Let us prove (S 2). Let $X \xrightarrow{f} Y \rightarrow Z' \rightarrow X[1]$ and $Y \xrightarrow{g} Z \rightarrow X' \rightarrow Y[1]$ be two distinguished triangles, with $X' \in \mathcal{N}$, $Z' \in \mathcal{N}$. By (TR 2) there exists a distinguished triangle $X \xrightarrow{g \circ f} Z \rightarrow Y' \rightarrow X[1]$, and by (TR 5) there exists a distinguished triangle $Z' \rightarrow Y' \rightarrow X' \rightarrow Z'[1]$. By (N 2), (N 3) and (TR 3), we have: $Y' \in \mathcal{N}$. Hence $g \circ f \in S(\mathcal{N})$.

To prove (S 3), consider a distinguished triangle $Z \xrightarrow{g} Y \xrightarrow{k} X' \rightarrow Z[1]$, with $X' \in \mathcal{N}$, and let $f : X \rightarrow Y$. There exists a distinguished triangle

$$W \xrightarrow{h} X \xrightarrow{k \circ f} X' \rightarrow W[1] .$$

Then by (TR 4) and (TR 3), we have a morphism of distinguished triangles:

$$\begin{array}{ccccccc} W & \xrightarrow{h} & X & \xrightarrow{k \circ f} & X' & \longrightarrow & W[1] \\ \downarrow & & \downarrow f & & \downarrow \text{id}_{X'} & & \downarrow \\ Z & \xrightarrow{g} & Y & \xrightarrow{k} & X' & \longrightarrow & Z[1] . \end{array}$$

Since $X' \in \mathcal{N}$, h belongs to $S(\mathcal{N})$.

A similar proof holds by reversing the arrows.

Finally we prove (S4). Let $f : X \rightarrow Y$ and $t : Y \rightarrow Y'$, with $t \in S(\mathcal{N})$ and $t \circ f = 0$. We shall show that there exists $s : X' \rightarrow X$, $s \in S(\mathcal{N})$, such that $f \circ s = 0$. Let $Z \xrightarrow{g} Y \xrightarrow{t} Y' \rightarrow Z[1]$ be a distinguished triangle, with $Z \in \mathcal{N}$. By (TR 1), (TR 3), (TR 4), there exists $h : X \rightarrow Z$ such that $f = g \circ h$. If we embed h into a distinguished triangle $X' \xrightarrow{s} X \xrightarrow{h} Z \rightarrow X'[1]$ then s will satisfy the desired properties. The proof of the converse implication is similar. \square

Notation 1.6.8. Let \mathcal{C} be a triangulated category and \mathcal{N} a null system in \mathcal{C} . We write \mathcal{C}/\mathcal{N} instead of $\mathcal{C}_{S(\mathcal{N})}$.

Proposition 1.6.9. *Let \mathcal{C} be a triangulated category and \mathcal{N} a null system.*

- (i) \mathcal{C}/\mathcal{N} becomes a triangulated category by taking for distinguished triangles those isomorphic to the image of a distinguished triangle in \mathcal{C} .
- (ii) Denote by Q the natural functor $\mathcal{C} \rightarrow \mathcal{C}/\mathcal{N}$. We have $Q(X) \simeq 0$ for $X \in \mathcal{N}$.
- (iii) Any functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ of triangulated categories such that $F(X) \simeq 0$ for all $X \in \mathcal{N}$, factors uniquely through Q .

Похідна категорія

$\mathcal{C} = \mathbb{k}\text{-mod}$.

We shall apply the preceding construction to the triangulated category $\mathbf{K}(\mathcal{C})$. It is clear that:

$$(1.7.1) \quad \mathcal{N} = \{X \in \text{Ob}(\mathbf{K}(\mathcal{C})); H^n(X) = 0 \text{ for any } n\}$$

is a null system. Note that, in view of Proposition 1.5.6, $S(\mathcal{N})$ consists of quasi-isomorphisms of $\mathbf{K}(\mathcal{C})$.

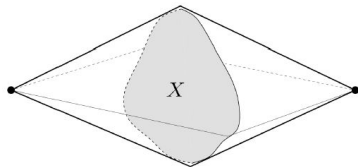
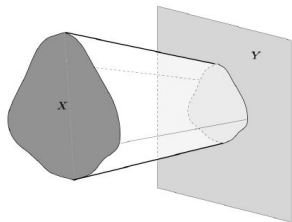
Definition 1.7.1. We set $\mathbf{D}(\mathcal{C}) = \mathbf{K}(\mathcal{C})/\mathcal{N}$ and call $\mathbf{D}(\mathcal{C})$ the derived category of \mathcal{C} .

By replacing $\mathbf{K}(\mathcal{C})$ with $\mathbf{K}^b(\mathcal{C})$ (resp. $\mathbf{K}^+(\mathcal{C})$, resp. $\mathbf{K}^-(\mathcal{C})$), we define similarly the derived categories $\mathbf{D}^b(\mathcal{C})$ (resp. $\mathbf{D}^+(\mathcal{C})$, resp. $\mathbf{D}^-(\mathcal{C})$). By Proposition 1.6.3 the functor $H^n(\cdot): \mathbf{K}(\mathcal{C}) \rightarrow \mathcal{C}$, factors through $\mathbf{D}(\mathcal{C})$. We still denote by $H^n(\cdot)$ the functor from $\mathbf{D}(\mathcal{C})$ to \mathcal{C} so obtained.

Proposition 1.7.2. (i) $\mathbf{D}^b(\mathcal{C})$ (resp. $\mathbf{D}^+(\mathcal{C})$, resp. $\mathbf{D}^-(\mathcal{C})$) is equivalent to the full subcategory of $\mathbf{D}(\mathcal{C})$ consisting of objects X such that $H^n(X) = 0$ for $|n| \gg 0$ (resp. $n \ll 0$, resp. $n \gg 0$).

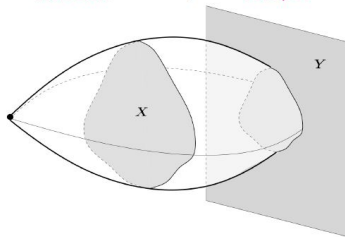
(ii) By the composition of the functors $\mathcal{C} \rightarrow \mathbf{K}(\mathcal{C}) \rightarrow \mathbf{D}(\mathcal{C})$, \mathcal{C} is equivalent to the full subcategory of $\mathbf{D}(\mathcal{C})$ consisting of objects X such that $H^n(X) = 0$ for $n \neq 0$.

Конус відображення в топології



$$X \xrightarrow{f} Y \xrightarrow[\text{inclusion}]{i} C_f \xrightarrow[\text{collapse}]{q} \Sigma X \text{ suspension}$$

mapping cone



Конус відображення як функтор

$\text{Cone} = M$ є функтором $\mathbf{dg}^{\rightarrow} \rightarrow \mathbf{dg}$, де $\rightarrow =$ категорія з двома об'єктами $0, 1$ і єдиною нетотожною стрілкою $0 \rightarrow 1$.

$\mathcal{C}^{\rightarrow} =$ категорія стрілок в \mathcal{C} :

$\text{Ob } \mathcal{C}^{\rightarrow} = \text{Mor } \mathcal{C}$,

$\mathcal{C}^{\rightarrow}(f : A \rightarrow B, g : X \rightarrow Y)$

$$= \left\{ (u : A \rightarrow X, v : B \rightarrow Y) \in (\text{Mor } \mathcal{C})^2 \left| \begin{array}{ccc} A & \xrightarrow{f} & B \\ u \downarrow & = & \downarrow v \\ X & \xrightarrow{g} & Y \end{array} \right. \right\}.$$

$$\text{Cone}(f : A \rightarrow B) = \left(A[1] \oplus B, \begin{pmatrix} d_{A[1]} & \sigma^{-1} \cdot f \\ 0 & d_B \end{pmatrix} \right).$$

Exercise

$$\begin{pmatrix} d_{A[1]} & \sigma^{-1} \cdot f \\ 0 & d_B \end{pmatrix}^2 = 0$$

$$\text{Cone} \left(\begin{array}{ccc} A & \xrightarrow{f} & B \\ u \downarrow & = & \downarrow v \\ X & \xrightarrow{g} & Y \end{array} \right) = u[1] \oplus v. \text{ Чому це } \in \mathbf{dg}?$$

Коротка точна послідовність комплексів індукує виділений трикутник в похідній категорії

Proposition 1.7.5. *Let \mathcal{C} be an abelian category and let $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0$ be an exact sequence in $\mathbf{C}(\mathcal{C})$. Let $M(f)$ be the mapping cone of f and let $\phi^n : M(f)^n = X^{n+1} \oplus Y^n \rightarrow Z^n$ be the morphism $(0, g^n)$. Then $\{\phi^n\}_n : M(f) \rightarrow Z$ is a morphism of complexes, $\phi \circ \alpha(f) = g$, and ϕ is a quasi-isomorphism.*

Proof. It is straightforward to see that ϕ is a morphism of complexes. Moreover we have an exact sequence:

$$0 \rightarrow M(\text{id}_X) \xrightarrow{\gamma} M(f) \rightarrow Z \rightarrow 0$$

where γ is associated to the morphism $\text{id}_X \rightarrow f$. This last morphism is described by the commutative diagram:

$$\begin{array}{ccc} X & \xrightarrow{\text{id}} & X \\ \downarrow \text{id} & & \downarrow f \\ X & \xrightarrow{f} & Y \end{array} .$$

By Proposition 1.3.6 it is enough to check that $H^n(M(\text{id}_X)) = 0$ for all $n \in \mathbb{Z}$. Since $M(\text{id}_X)$ is zero in $\mathbf{K}(\mathcal{C})$, this is evident. \square

$$\begin{array}{ccc}
 X & \xrightarrow{1} & X \\
 \downarrow 1 & & \downarrow f \\
 X & \xrightarrow{f} & Y
 \end{array}
 \xrightarrow{\text{Cone}} \text{Cone}(1, f) = 1 \oplus f : \text{Cone}1_X \rightarrow \text{Cone}f.$$

$$\exists \phi = \begin{pmatrix} 0 \\ g \end{pmatrix} : \text{Cone}f \rightarrow Z \in \mathbf{dg} \Leftrightarrow \begin{pmatrix} d_{X[1]} & \sigma^{-1} \cdot f \\ 0 & d_Y \end{pmatrix} \begin{pmatrix} 0 \\ g \end{pmatrix} = \begin{pmatrix} 0 \\ g \end{pmatrix} d_Z.$$

$$0 \rightarrow \text{Cone}1_X \xrightarrow{\text{Cone}(1, f)} \text{Cone}f \xrightarrow{\phi} Z \rightarrow 0$$

точна як \oplus в \mathbf{gr} послідовностей

$$0 \rightarrow X[1] \xrightarrow{1} X[1] \rightarrow 0 \text{ і } 0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0.$$

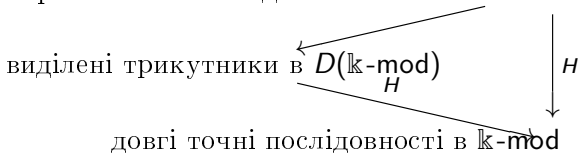
Exercise

$\text{Cone}1_X$ – стягуваний ($\exists s : 1_{\text{Cone}1} = sd + ds$).

Тому $\text{Cone}1_X$ – ациклічний $\Rightarrow \phi$ – квазі-ізоморфізм $\Rightarrow \phi$ оборотний в $D(\mathbb{k}\text{-mod}) \Rightarrow$ в $D(\mathbb{k}\text{-mod})$ існує ізоморфізм трикутників

$$\begin{array}{ccccccc}
 X & \xrightarrow{f} & Y & \xrightarrow{g} & Z & \xrightarrow{\phi^{-1} \cdot \beta(f)} & X[1] \\
 \parallel & & \parallel & & \uparrow \cong & & \parallel \\
 X & \xrightarrow{f} & Y & \xrightarrow{\alpha(f)} & \text{Cone } f & \xrightarrow{\beta(f)} & X[1]
 \end{array}$$

короткі точні послідовності комплексів \mathbb{k} -модулів



In the situation of Proposition 1.7.5 the distinguished triangle $X \rightarrow Y \rightarrow Z \xrightarrow{h} X[1]$ is called the distinguished triangle associated to the exact sequence $0 \rightarrow X \rightarrow Y \rightarrow Z \rightarrow 0$. Here $h = \beta(f) \circ \phi^{-1}$.

Note that the above distinguished triangle gives rise to a long exact sequence:

$$\cdots \longrightarrow H^n(X) \longrightarrow H^n(Y) \longrightarrow H^n(Z) \xrightarrow{H^n(h)} H^{n+1}(X) \longrightarrow \cdots$$

and $H^n(h) = -\delta$, δ being defined in Proposition 1.3.6.



Masaki Kashiwara and Pierre Schapira, Sheaves on manifolds, Grundlehren der mathematischen Wissenschaften, vol. 292, Springer-Verlag, Berlin, New York, 1990. §1.4.4–§1.7.5